

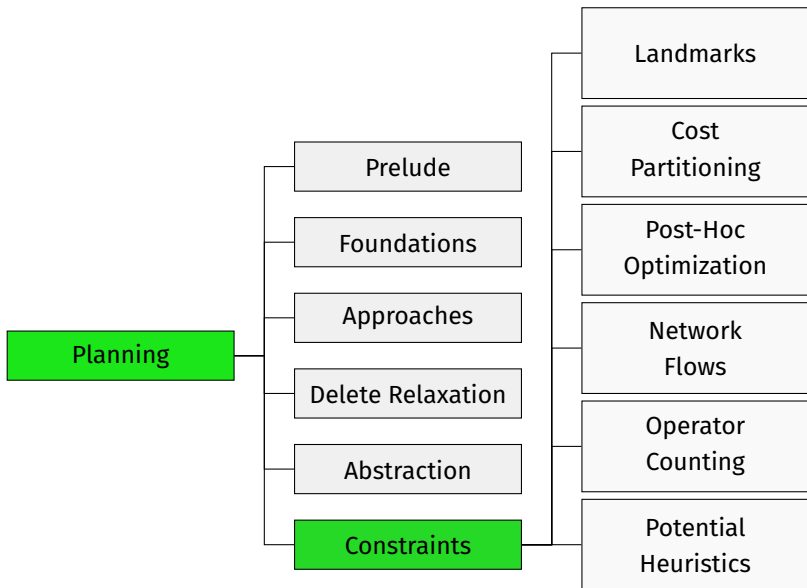
# Automated Planning

## F1. Constraints: Introduction

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# Content of this Course



# Constraint-based Heuristics

## Coming Up with Heuristics in a Principled Way

### General Procedure for Obtaining a Heuristic

Solve a simplified version of the problem.

Major ideas for heuristics in the planning literature:

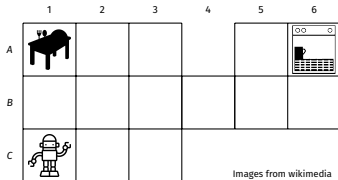
- delete relaxation
- abstraction
- critical paths
- landmarks
- network flows
- potential heuristic

Landmarks, network flows and potential heuristics are based on **constraints** that can be specified for a planning task.

## Constraints: Example

FDR planning task  $\langle V, I, O, \gamma \rangle$  with

- $V = \{robot-at, dishes-at\}$  with
  - $dom(robot-at) = \{A1, \dots, C3, B4, A5, \dots, B6\}$
  - $dom(dishes-at) = \{Table, Robot, Dishwasher\}$
- $I = \{robot-at \mapsto C1, dishes-at \mapsto Table\}$
- operators
  - move-x-y to move from cell x to adjacent cell y
  - pickup dishes, and
  - load dishes into the dishwasher.
- $\gamma = (robot-at = B6) \wedge (dishes-at = Dishwasher)$



Images from wikimedia

# Constraints

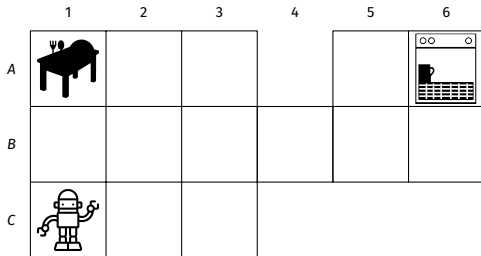
Some heuristics exploit **constraints** that describe something that holds in every solution of the task.

For instance, every solution is such that

- a variable takes a certain value in at least one visited state.  
(a **fact landmark** constraint)

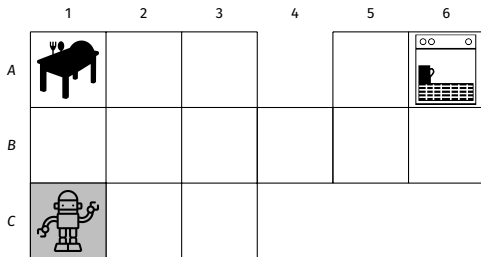
# Fact Landmarks: Example

Which values do *robot-at* and *dishes-at* take in every solution?



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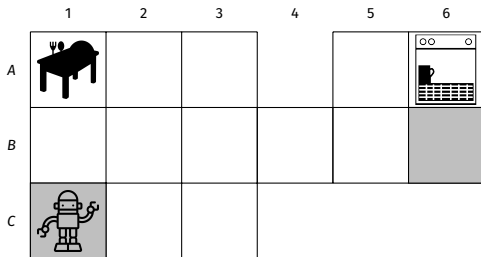


- $\text{robot-at} = C1, \text{dishes-at} = \text{Table}$  (initial state)



## Fact Landmarks: Example

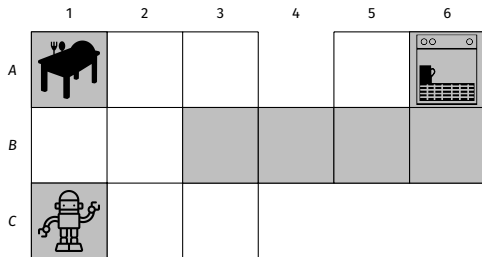
Which values do *robot-at* and *dishes-at* take in every solution?



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## Fact Landmarks: Example

Which values do *robot-at* and *dishes-at* take in every solution?



- *robot-at* = C1, *dishes-at* = Table (initial state)
- *robot-at* = B6, *dishes-at* = Dishwasher (goal state)
- *robot-at* = A1, *robot-at* = B3, *robot-at* = B4, *robot-at* = B5, *robot-at* = A6, *dishes-at* = Robot

# Constraints

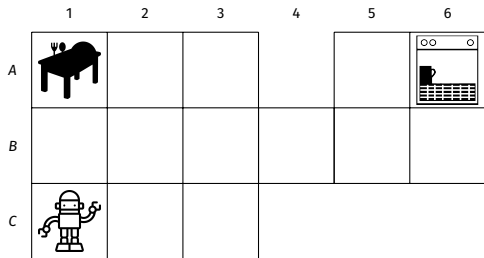
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- an action must be applied.  
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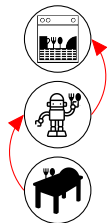
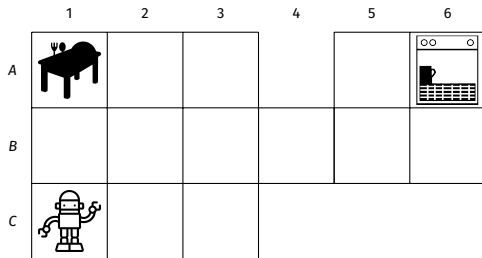
# Action Landmarks: Example

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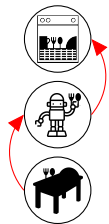
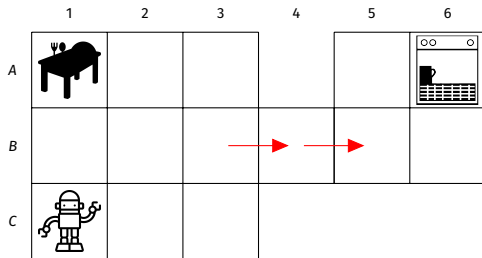
Which actions must be applied in every solution?



- pickup
- load

# Action Landmarks: Example

Which actions must be applied in every solution?



- pickup
- load
- move-B3-B4
- move-B4-B5

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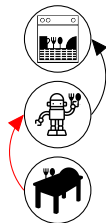
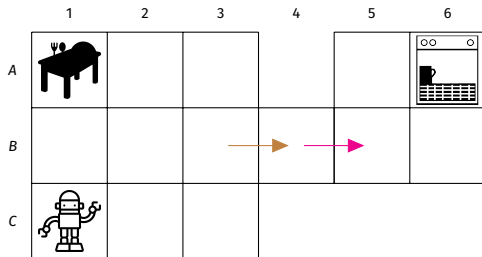
For instance, every solution is such that

- a variable takes some **value** in at least one visited state.  
(a **fact landmark** constraint)
- at least one action from a set of actions must be applied.  
(a **disjunctive action landmark** constraint)



## Disjunctive Action Landmarks: Example

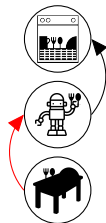
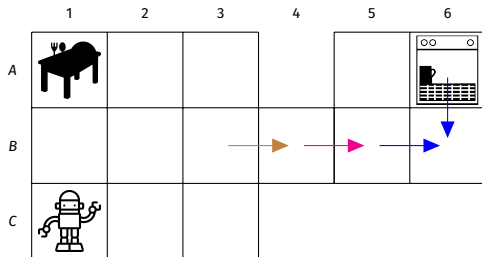
Which set of actions is such that at least one must be applied?



- {pickup}
- {load}
- {move-B3-B4}
- {move-B4-B5}

# Disjunctive Action Landmarks: Example

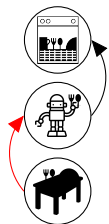
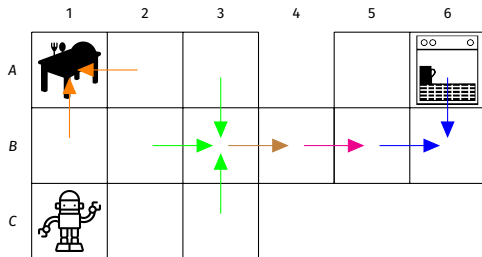
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- {pickup}
- {load}
- {move-B3-B4}
- {move-B4-B5}
- {move-A6-B6, move-B5-B6}

# Disjunctive Action Landmarks: Example

Which set of actions is such that at least one must be applied?



- {pickup}
- {load}
- {move-B3-B4}
- {move-B4-B5}
- {move-A6-B6, move-B5-B6}
- {move-A3-B3, move-B2-B3, move-C3-B3}
- {move-B1-A1, move-A2-A1}
- ...

# Constraints

Some heuristics exploit **constraints** that describe something that holds in every solution of the task.

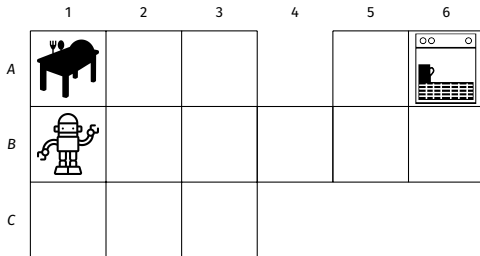
For instance, every solution is such that

- a variable takes some value in at least one visited state.  
(a **fact landmark** constraint)
- at least one action from a set of actions must be applied.  
(a **disjunctive action landmark** constraint)
- fact consumption and production is “balanced”.  
(a **network flow** constraint)

## Network Flow: Example

Consider the fact robot-at = B2.

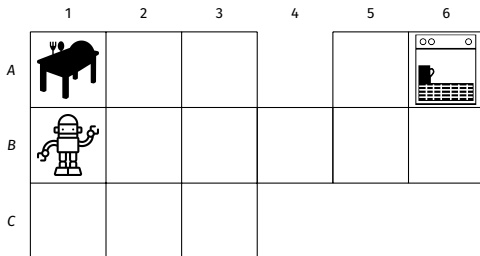
How often are actions used that enter this cell?



## Network Flow: Example

Consider the fact  $\text{robot-at} = B2$ .

How often are actions used that enter this cell?



**Answer:** as often as actions that leave this cell

If  $\text{Count}_o$  denotes how often operator  $o$  is applied, we have:

$$\text{Count}_{\text{move-A1-B1}} + \text{Count}_{\text{move-B2-B1}} + \text{Count}_{\text{move-C1-B1}} =$$

$$\text{Count}_{\text{move-B1-A1}} + \text{Count}_{\text{move-B1-B2}} + \text{Count}_{\text{move-B1-C1}}$$

# Multiple Heuristics

# Combining Admissible Heuristics Admissibly

Major ideas to combine heuristics admissibly:

- maximize
- canonical heuristic (for abstractions)
- **minimum hitting set** (for landmarks)
- **cost partitioning**
- **operator counting**

Often computed as solution to a **(integer) linear program**.



## Combining Heuristics Admissibly: Example

### Example

Consider an FDR planning task  $\langle V, I, \{o_1, o_2, o_3, o_4\}, \gamma \rangle$  with  $V = \{v_1, v_2, v_3\}$  with  $\text{dom}(v_1) = \{A, B\}$  and  $\text{dom}(v_2) = \text{dom}(v_3) = \{A, B, C\}$ ,  $I = \{v_1 \mapsto A, v_2 \mapsto A, v_3 \mapsto A\}$ ,

$$o_1 = \langle v_1 = A, v_1 := B, 1 \rangle$$

$$o_2 = \langle v_2 = A \wedge v_3 = A, v_2 := B \wedge v_3 := B, 1 \rangle$$

$$o_3 = \langle v_2 = B, v_2 := C, 1 \rangle$$

$$o_4 = \langle v_3 = B, v_3 := C, 1 \rangle$$

and  $\gamma = (v_1 = B) \wedge (v_2 = C) \wedge (v_3 = C)$ .

Consider all atomic projections. By using additivity for orthogonal abstractions, which heuristic estimates can we sum up admissibly?

## Combining Heuristics Admissibly: Example

### Example

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and  $\gamma = (v_1 = B) \wedge (v_2 = C) \wedge (v_3 = C)$ .

Consider all atomic projections. By using additivity for orthogonal abstractions, which heuristic estimates can we sum up admissibly?

**Answer:** Let  $h_i := h^{v_i}$ . Then  $h = \max \{h_1 + h_2, h_1 + h_3\}$  is admissible.

## Reminder: Orthogonality and Additivity

Why can we add  $h_1$  and  $h_2$  ( $h_1$  and  $h_3$ ) admissibly?

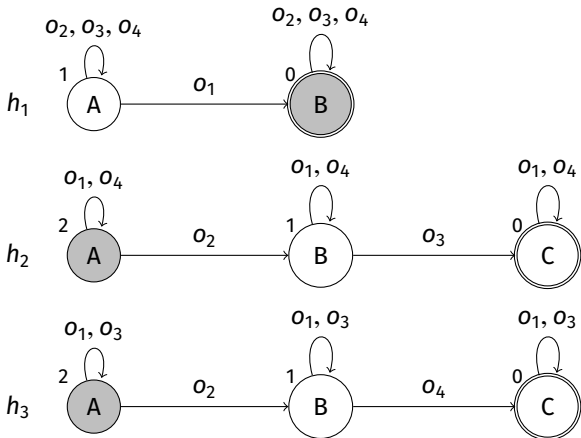
### Theorem (Additivity for Orthogonal Abstractions)

*Let  $h^{\alpha_1}, \dots, h^{\alpha_n}$  be abstraction heuristics of the same transition system such that  $\alpha_i$  and  $\alpha_j$  are orthogonal for all  $i \neq j$ .*

*Then  $\sum_{i=1}^n h^{\alpha_i}$  is a safe, goal-aware, admissible and consistent heuristic for  $\Pi$ .*

# Combining Heuristics (In)admissibly: Example

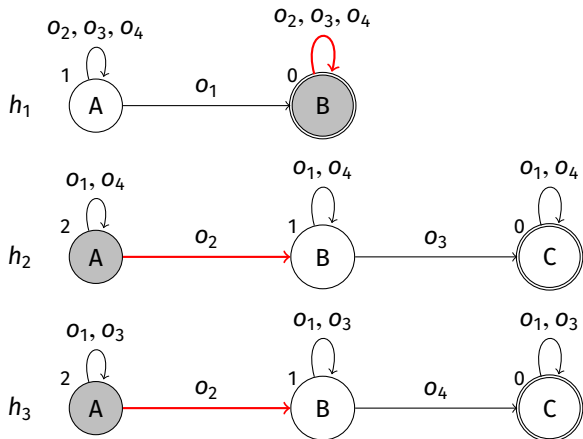
Let  $h = h_1 + h_2 + h_3$ .



$\langle o_2, o_3, o_4 \rangle$  is a plan for  $s = \langle B, A, A \rangle$  but  $h(s) = 4$ .

# Combining Heuristics (In)admissibly: Example

Let  $h = h_1 + h_2 + h_3$ .



$\langle o_2, o_3, o_4 \rangle$  is a plan for  $s = \langle B, A, A \rangle$  but  $h(s) = 4$ .

Heuristics  $h_2$  and  $h_3$  both account for the single application of  $o_2$ .

## Prevent Inadmissibility

The reason that  $h_2$  and  $h_3$  are not additive is because the cost of  $o_2$  is considered in both.

Is there anything we can do about this?

## Prevent Inadmissibility

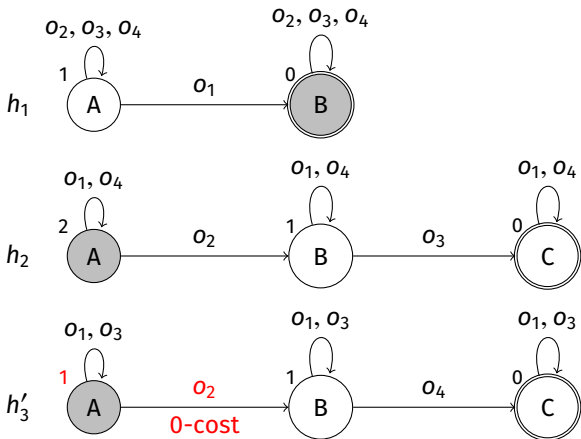
The reason that  $h_2$  and  $h_3$  are not additive is because the cost of  $o_2$  is considered in both.

Is there anything we can do about this?

**Solution:** We can ignore the cost of  $o_2$  in one heuristic by setting its cost to 0 (e.g.,  $cost_3(o_2) = 0$ ).

## Combining Heuristics Admissibly: Example

Let  $h' = h_1 + h_2 + h'_3$ , where  $h'_3 = h^{V_3}$  assuming  $cost_3(o_2) = 0$ .



$\langle o_2, o_3, o_4 \rangle$  is an optimal plan for  $s = \langle B, A, A \rangle$  and  $h'(s) = 3$  an admissible estimate.



## Cost partitioning

Using the cost of every operator only in one heuristic is called a **zero-one cost partitioning**.

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More generally, heuristics are additive if all operator costs are distributed in a way that the sum of the individual costs is no larger than the cost of the operator.

This can also be expressed as a constraint, the **cost partitioning constraint**:

$$\sum_{i=1}^n cost_i(o) \leq cost(o) \text{ for all } o \in O$$

(more details later)

# Summary

## Summary

- Landmarks and network flows are **constraints** that describe something that holds in every solution of the task.
- Heuristics can be combined admissibly if the **cost partitioning constraint** is satisfied.